

# Nonuniform Traffic Distribution Model in Reverse Link of Multirate/Multiservice WCDMA-Based Systems

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**Abstract**—This paper focuses on the modeling of the reverse link of a wideband code division multiple access system in a nonhomogeneous environment with a single cell. Multiple traffic spatial and service nonuniformities are considered in the analytical model, and then, expressions for required transmitted power and the associated outage probability and block error rate are derived. Special attention is also paid to the effect caused by different transmission bit rates and the spatial location of the traffic nonuniformities. From the presented expressions, it is possible to set appropriate load thresholds to control the desired error rate. Although the model considers a single cell, results in terms of maximum allowable load can also be applicable in multicell scenarios.

**Index Terms**—Load factor, nonuniformities, service distribution, traffic distribution, wideband code division multiple access (WCDMA) capacity.

## I. INTRODUCTION

ONE of the main traffic characteristics in cellular networks is nonhomogeneous traffic spatial distribution. To cope with this inherent aspect in mobile communication scenarios, system deployment must be preceded by careful network planning. Network planning takes into account the expected traffic load all over the service area in order to come up with a solution on how many base stations are required and what are the suitable locations to set up the sites. Certain operator's objectives in coverage, capacity, and quality must be met in the planning exercise. Furthermore, in the context of multimedia services, traffic needs to be detailed also in terms of traffic class spatial distribution (e.g., expected conversational traffic and interactive traffic) because of the different quality of service (QoS) requirements.

Nevertheless, although network planning can consider these nonhomogeneities at some extent, the high dynamics associated with the traffic generation processes coupled with user mobility clearly need additional mechanisms to cope with the potential problems with network performance derived from traffic profile distributions that are significantly different from

those expected in the network planning phase. Traffic level and service mix variations along time and space, thus resulting in different load levels in different cells and times, may result in an observed QoS that is significantly different from the planned values, hence negatively impacting the user's mobile experience.

In the context of wideband code division multiple access (WCDMA) systems, service and traffic spatial distributions become even more important, since there is a tight dependence of the interference level observed on every user's position in the network and the corresponding transmitted power levels. Clearly, the research community has already targeted these aspects. First approaches may be found in [1] and [2], where the impact of certain traffic distributions is analyzed by simulations and analytical expressions. This paper is focused on the characterization of the effect of nonuniformities in the reverse link of a WCDMA system. Thus, the major difference with respect to [1] and [2] lies in the fact that in these papers, no service mix is considered, and therefore, only specific traffic spatial distributions are analyzed. On the contrary, this paper presents an accurate traffic description that allows the specification of a generic scenario by breaking it down into different traffic layers, and therefore, general expressions for spatial distributions are obtained. Moreover, each layer is associated with a specific and independent service distribution. Required power expressions that have been obtained are general, and thus, they can be applied to any scenario regardless of the number of layers it consists of. This paper is an evolution of a previous work [3], where simulations showed that nonuniformities meant a new challenge to be faced in terms of capacity and a scenario with a single hotspot was analyzed by simulations. In this context, the formulation of the general problem is developed throughout this paper.

Other approaches presented in the literature also cope with techniques to mitigate the effects caused by nonuniformities. Radio resource management (RRM) strategies have become one of the targets [4], specifically call admission control (CAC) and load control, on which many studies have been carried out. In [5] and [6], different algorithms are presented and simulated, achieving positive results in maximizing the capacity while minimizing parameters such as blocking probability. Finally, pilot power management has also emerged as a solution to balance the load among base stations in nonhomogeneous scenarios. For instance, in [7]–[11], different algorithms are proposed, and all of them focused on forward link. In this respect,

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this paper may also serve as the basis for further development of these RRM techniques (e.g., devising admission and congestion thresholds), as long as the analytical approach allows us to gain insight into the general problem. RRM and CAC, in particular, are of crucial importance for the correct behavior of the system. The limitation of the interference lies on this mechanism and allows guaranteeing a maximum load level that should not be exceeded. These load thresholds must be determined by taking into account aspects such as the service used by users, their activity factor, and their spatial distribution. In fact, their huge variability is shown in this paper.

As aforementioned, this paper focuses on modeling traffic nonhomogeneities from an analytical perspective, so that the key issues affecting the problem can be extracted. The final objective is to propose a model able to compute the maximum capacity (i.e., maximum uplink load factor) that can be supported in the uplink of a WCDMA cell under nonhomogeneous spatial traffic distribution and multiple services. This maximum load factor value can then be used as an input to dynamic RRM strategies, e.g., admission or load control. Capacity will be defined for specific outage probability requirements. The analytical expressions have been derived for a single cell model. However, as it will be discussed in this paper, results in terms of maximum load factor could also be applicable in a multicell scenario. Additionally, the envisaged impact of soft handover is discussed. The remainder of this paper is organized as follows: The problem formulation is presented in Section II. Section III proposes a model for computing the transmitted power distribution in nonuniformly distributed scenarios and derives outage probability. Finally, Section IV presents results for some representative case studies, and Section V presents the conclusions.

## II. PROBLEM FORMULATION

The considered scenario is an isolated WCDMA cell with radius  $R$  (in meters). This scenario is composed of  $N$  overlapped layers ( $L_i$ , with  $0 \leq i \leq N - 1$ ), each one with specific traffic and service characteristics. All layers are circular shaped with a radius  $r_i$  (in meters) and their center is placed  $D_i$  (in meters) from the cell site. In particular, layer  $L_0$  matches up with the cell area, and thus,  $D_0 = 0$  and  $r_0 = R$  (the cell site is set as the coordinate origin). The scenario layout is shown in Fig. 1.

Let us denote  $T_T$  as the total number of users within the cell.  $T_T$  is distributed among the  $N$  layers, and the proportion of users belonging to layer  $L_i$  is given by

$$\alpha_i = \frac{T_i}{\sum_{j=0}^{K-1} T_j} = \frac{T_i}{T_T} \quad (1)$$

where  $T_i$  is the number of users within layer  $L_i$ . On the other hand,  $K$  possible services ( $0 \leq n \leq K - 1$ ) are identified.  $\rho_{i,n}$  denotes the proportion of users with the  $n$ th service in layer  $L_i$ . In general, the traffic source is not continuously transmitting packets within a session, and some activity periods alternate with inactivity periods (e.g., reading time during a World Wide Web download or silence periods in speech calls). Then, the

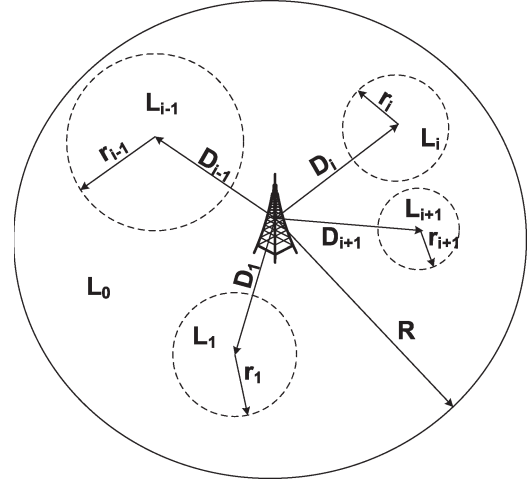


Fig. 1. Scenario layout.

activity factor for the  $n$ th service ( $\epsilon_n$ ) is defined as the proportion of time in which a user is transmitting for this service, taking into account session and intersession times.

With respect to activity periods in which a certain user is transmitting data through the air interface, there may exist several possibilities for each service to transmit the data flow, which is characterized essentially by a transmission bit rate, a channel code, and a required  $E_b/N_0$ . Following the Third Generation Partnership Project (3GPP) terminology [13], these possibilities will be denoted as a transport format (TF, hereafter).<sup>1</sup> Each TF is characterized by a transmission bit rate and an  $E_b/N_0$  target. Then, the set of TFs is denoted as a TF set (TFS). In particular,  $TFS_n$  stands for the TFS of the  $n$ th service. The set of possible bit rates for the  $n$ th service is  $R_{b_{n,j}}$  ( $0 \leq j \leq TFS_n - 1$ ), and the corresponding  $E_b/N_0$  is  $(E_b/N_0)_{n,j}$  with  $0 \leq j \leq TFS_n - 1$ .  $p_{n,j}$  is defined as the probability that the  $n$ th service uses the  $j$ th TF, which depends on the specific medium access control (MAC) algorithm that executes the TF selection in the uplink (see, e.g., [4] with some examples).

In Section III, an analytical model of the reverse link of a multirate/multiservice WCDMA-based system is developed. The main target of this paper is to obtain a model to study the impact of nonuniformities on the performance of the system. The QoS parameters under study are the outage probability, which is defined as the probability that the measured  $(E_b/N_0)_m$  is below the target, and the block error rate (BLER), which is defined as the percentage of erroneous blocks received by the base station. Notice that outage and BLER are tightly coupled with transmitted power, and thus, the study of the transmitted power becomes the key issue.

The process carried out in Sections III-A and B is aimed at determining the required transmitted power probability density function (pdf) expression to evaluate the outage probability and BLER performance.

<sup>1</sup>Notice that, although the 3GPP terminology is used here, the concept of TF would be general to any code division multiple access system with variable transmission bit rate.

### III. ANALYTICAL MODEL

#### A. Transmitted Power Distribution

In WCDMA-based systems, the capacity is highly dependent on the interference existing in the air interface. The main goal of this section is to determine the expression of the required transmitted power pdf to derive the system capacity. In the uplink case, the required transmitted power (in logarithmic units) is given by [4]

$$P(\text{dBm}) = Z(\text{dB}) + 10 \log \left( \frac{P_N}{(1-\eta) \left( \frac{W}{\left( \frac{E_b}{N_0} \right) R_b} + 1 \right)} \right) \\ = Z + \Phi \quad (2)$$

where  $W$  is the total bandwidth after spreading,  $P_N$  is the thermal noise power,  $Z$  is the total propagation loss (path loss plus a shadowing component),  $E_b/N_0$  is the target energy per bit versus interference plus noise spectral density, and  $\eta$  is the uplink load factor [3]. The load factor is defined as the proportion of the power generated by users (own-cell power and other-cells interference) with respect to the total power (thermal noise power, own-cell received power, and other-cell interference), and thus, it is tightly coupled with the maximum number of users that can be allocated in the cell (the maximum cell capacity), i.e.,

$$\eta = \frac{P_R + P_I}{P_N + P_R + P_I} \\ = 1 - \frac{P_N}{P_N + P_R + P_I} \quad (3)$$

where  $P_R$  is the total own-cell received power, and  $P_I$  is the total other-cell interference. Particularizing (3) for a single cell scenario, the load factor is given by

$$\eta = \frac{P_R}{P_N + P_R} \\ = 1 - \frac{P_N}{P_N + P_R}. \quad (4)$$

In general, RRM algorithms in the uplink tend to control the load factor. Thus, the presented model intends to determine the maximum admissible load factor on the basis of traffic distribution and desired quality.

For a given load factor  $\eta$ , the required transmitted power becomes a function of total loss,  $E_b/N_0$ , and  $R_b$ . Thus, the required transmitted power ( $P$ ) pdf may be obtained by calculating the pdf of the total loss  $Z$ , which depends on traffic spatial distribution, and of the product  $(E_b/N_0) \cdot R_b$ , which depends on service characterization. The total propagation loss is defined as

$$Z(\text{dB}) = Y(\text{dB}) + S(\text{dB}) \quad (5)$$

where  $Y$  is the path loss dependent on the distance, and  $S$  is the lognormal shadowing component (assuming that closed-loop power control cancels all fast fading effects). The path loss at

distance  $d$  (in meters) from the transmitter in logarithmic units  $Y$  (in decibels) is typically given by

$$Y(\text{dB}) = Y_0 + \zeta \log d \quad (6)$$

where  $Y_0$  and  $\zeta$  are constants and depend on the environment. Then, the path loss distribution is determined by the traffic spatial distribution (i.e., the distribution of the distance  $d$ ). Focusing on a generic layer  $L_i$  (characterized by its radius  $r_i$  and the distance  $D_i$  from the cell site to its center), the path loss pdf, which is denoted as  $f_Y^i(y)$ , is given as follows (see the Appendix for details):

if  $D_i \geq r_i$ :

$$f_Y^i(y) = \frac{A\beta}{2\pi r_i^2} e^{\beta y} \left[ \pi - 2 \arcsin \frac{D_i^2 - r_i^2 + Ae^{\beta y}}{2D_i \sqrt{(A)e^{\beta y}}} \right] \quad (7)$$

in the range  $Y_0 + \zeta \log(D_i - r_i) \leq y \leq Y_0 + \zeta \log(D_i + r_i)$ ;

if  $D_i < r_i$ :

$$f_Y^i(y) = \frac{A\beta}{r_i^2} e^{\beta y} \quad (8)$$

in the range  $-\infty < y \leq Y_0 + \zeta \log(r_i - D_i)$ ; and

$$f_Y^i(y) = \frac{A\beta}{2\pi r_i^2} e^{\beta y} \left[ \pi - 2 \arcsin \frac{D_i^2 - r_i^2 + Ae^{\beta y}}{2D_i \sqrt{(A)e^{\beta y}}} \right] \quad (9)$$

in the range  $Y_0 + \zeta \log(r_i - D_i) < y \leq Y_0 + \zeta \log(D_i + r_i)$ , where  $A = 10^{-2(Y_0/\zeta)}$ , and  $\beta = 2(\ln(10)/\zeta)$ .

Adding the effects of lognormal shadowing  $S$ , the total propagation loss ( $Z = Y + S$ ) pdf is the result of convolution (7)–(9) and a Gaussian function with null mean and  $\sigma^2$  variance, so that the resulting pdf can be expressed as

$$f_Z^i(z) = f_Y^i(z) * \frac{1}{\sqrt{2\pi\sigma}} e^{-\frac{z^2}{2\sigma^2}} \quad (10)$$

where  $*$  denotes the convolutional operator.

Finally, the overall total propagation loss density function is obtained by summing each pdf that is weighted according to the traffic fraction  $\alpha_i$  of each layer. In particular, when  $N$  layers are considered (from 0 to  $N - 1$ ), the final total propagation loss density function is given by

$$f_Z(z) = \sum_{i=0}^{N-1} \alpha_i f_Z^i(z). \quad (11)$$

It may be observed from (2) that the required transmitted power is also dependent on  $\Phi$ . In the case of WCDMA-based systems, only transmitting users (i.e., during activity periods) contribute to an increase of the interference. Taking this fact into account, the influence of each user on the total interference depends on the activity factor  $\epsilon_n$ , the services, the selected TF, and the distribution of services within the layers  $(\rho_{i,n})$ .

Therefore,  $\Phi$  is a random variable whose probability function is given by

$$P_{\Phi}(\Phi = \phi_{n,j}) = \sum_{i=0}^{N-1} P_{\Phi}^i(\phi_{n,j}) \quad (12)$$

where  $P_{\Phi}^i$  is given as follows:

$$P_{\Phi}^i(\phi_{n,j}) = \frac{\beta_n \alpha_i \rho_{i,n} p_{n,j}}{\sum_{s=0}^{K-1} \sum_{q=0}^{N-1} \alpha_q \beta_s \rho_{q,s}} \quad (13)$$

and  $\phi_{n,j}$  is defined as

$$\phi_{n,j} = 10 \log \left( \frac{P_N}{(1-\eta) \left( \frac{W}{\left( \frac{E_b}{N_0} \right)_{n,j} R_{b_{n,j}}} + 1 \right)} \right) \quad (14)$$

for  $j = 0, \dots, \text{TFS}_n - 1$ . Thus, (13) is defined in a discrete probability space  $\Omega_{\Phi} = \{\phi_{0,0}, \phi_{0,1}, \dots, \phi_{K-1}, \text{TFS}_{K-1-1}\}$ .

Finally, from (2), for a given load factor the required transmitted power pdf is obtained by convolutioning  $Z$  and  $\Phi$  pdfs, i.e.,

$$f_P(P) = f_Z(P) * f_{\Phi}(P). \quad (15)$$

Therefore, by combining (11) and (13), the required transmitted power density function results in

$$f_P(P) = \frac{\sum_{n=0}^{K-1} \sum_{j=0}^{\text{TFS}_n-1} \sum_{i=0}^{N-1} \alpha_i \epsilon_n \rho_{i,n} p_{n,j} f_Z^i(P - \phi_{n,j})}{\sum_{s=0}^{K-1} \sum_{q=0}^{N-1} \alpha_q \epsilon_s \rho_{q,s}}. \quad (16)$$

The required transmitted power pdf may be obtained by the following expressions for each layer and service, respectively:

$$f_P^i(P) = \frac{\sum_{n=0}^{K-1} \sum_{j=0}^{\text{TFS}_n-1} \epsilon_n \rho_{i,n} p_{n,j} f_Z^i(P - \phi_{n,j})}{\sum_{s=0}^{K-1} \epsilon_s \rho_{i,s}} \quad (17)$$

$$f_{P_n}(P) = \frac{\sum_{i=0}^{N-1} \sum_{j=0}^{\text{TFS}_n-1} \alpha_i \rho_{i,n} p_{n,j} f_Z^i(P - \phi_{n,j})}{\sum_{i=0}^{N-1} \alpha_i \rho_{i,n}}. \quad (18)$$

Notice that no power limitations have been taken into account so far. In real user equipment, transmitted power in the reverse link will be limited by an upper threshold ( $P_{\text{Tmax}}$ ) and a lower threshold ( $P_{\text{Tmin}}$ ). Then, the modified transmitted power

distribution taking into account power constraints ( $f_P^*(P)$ ) will be given by

$$f_P^*(P) = \begin{cases} 0, & P < P_{\text{Tmin}} \\ \int_{-\infty}^{P_{\text{Tmin}}} f_P(P) dP, & P = P_{\text{Tmin}} \\ f_P(P), & P_{\text{Tmin}} < P < P_{\text{Tmax}} \\ \int_{P_{\text{Tmax}}}^{\infty} f_P(P) dP, & P = P_{\text{Tmax}} \\ 0, & P > P_{\text{Tmax}} \end{cases} \quad (19)$$

## B. QoS Evaluation

This section extracts from the previous model the QoS performance in terms of outage probability  $\mu$  and BLER. It is assumed as an ideal power control so that the transmitted power is the required one, provided that it is between the limit  $[P_{\text{Tmin}}, P_{\text{Tmax}}]$ . A user is said to be in outage whenever the measured  $(E_b/N_0)_m$  is below the target value  $E_b/N_0$  or equivalently when  $P$  is above  $P_{\text{Tmax}}$ . Then, the outage probability  $\mu$  may be expressed as

$$\begin{aligned} \mu &= \text{prob} \left( \left( \frac{E_b}{N_0} \right)_m < \left( \frac{E_b}{N_0} \right) \right) \\ &= \text{prob}(P > P_{\text{Tmax}}). \end{aligned} \quad (20)$$

For a given load factor, the outage probability expression is given as follows:

$$\mu_{\eta} = \int_{P_{\text{Tmax}}}^{\infty} f_P(P) dP \quad (21)$$

$$\mu_{\eta} = \frac{\sum_{n=0}^{K-1} \sum_{j=0}^{\text{TFS}_n-1} \sum_{i=0}^{N-1} \alpha_i \epsilon_n \rho_{i,n} p_{n,j} \lambda_{i,n,j}}{\sum_{s=0}^{K-1} \epsilon_s \sum_{q=0}^{N-1} \alpha_q \rho_{q,s}} \quad (22)$$

with

$$\lambda_{i,n,j} = \int_{P_{\text{Tmax}}}^{\infty} f_Z^i(P - \phi_{n,j}) dP. \quad (23)$$

Likewise, BLER depicts the percentage of erroneous transport blocks (TB) received by the base station. TB is the minimum amount of information that can be exchanged between the physical and MAC layers. It includes a data part and the MAC and radio link control headers. BLER, which is given in the following expression, is a function of  $(E_b/N_0)_m$  and depends on physical layer aspects such as channel coding, modulation, and interleaving:

$$\text{BLER} = b_{n,j} \left( \left( \frac{E_b}{N_0} \right)_m \right) \quad (24)$$

where the function  $b_{n,j}(\cdot)$  is specific for each service and TF.

Whenever the required power  $P$  is between  $P_{\text{Tmin}}$  and  $P_{\text{Tmax}}$ , the user will be able to reach the target  $E_b/N_0$ , so that  $E_b/N_0 = (E_b/N_0)_m$ . On the contrary, when the required

power  $P$  is above the maximum available  $P_{T\max}$ , the measured  $(E_b/N_0)_m$  is lower than the target  $E_b/N_0$ . In particular,  $(E_b/N_0)_m$  will be  $(P - P_{T\max})$  dB below the target  $E_b/N_0$ . On the contrary, when  $P$  is below  $P_{T\min}$ , the difference between  $(E_b/N_0)_m$  and  $E_b/N_0$  is equal to  $P_{T\min} - P$ , and in this case, the measured  $(E_b/N_0)_m$  is above the target.

Summarizing these situations, the measured  $(E_b/N_0)_m$  as a function of the required transmitted power  $P$  is given by

$$\left(\frac{E_b}{N_0}\right)_m = g(P) = \begin{cases} \left(\frac{E_b}{N_0}\right) \frac{P_{T\min}}{P}, & P < P_{T\min} \\ \left(\frac{E_b}{N_0}\right), & P_{T\min} \leq P \leq P_{T\max} \\ \left(\frac{E_b}{N_0}\right) \frac{P_{T\max}}{P}, & P > P_{T\max} \end{cases} \quad (25)$$

Therefore, BLER is a function of the required transmitted power, i.e.,

$$\text{BLER} = \gamma_{n,j}(g(P)) = \gamma_{n,j}(P) \quad (26)$$

where  $\gamma_{n,j}(\cdot)$  is the function relating BLER and transmitted power for service  $n$  and TF  $j$ . In particular, for a given load factor, BLER is given as follows:

$$\begin{aligned} \overline{\text{BLER}}_\eta &= \int_{-\infty}^{\infty} \gamma(P) f_P(P) dP \\ \overline{\text{BLER}}_\eta &= \frac{\sum_{n=0}^{K-1} \sum_{j=0}^{TFS_n-1} \sum_{i=0}^{N-1} \alpha_i \epsilon_n \rho_{i,n} p_{n,j} \vartheta_{i,n,j}}{\sum_{s=0}^{K-1} \epsilon_s \sum_{q=0}^{N-1} \alpha_q \rho_{q,s}} \end{aligned} \quad (27)$$

with

$$\vartheta_{i,n,j} = \int_{-\infty}^{\infty} \gamma_{n,j}(P) f_Z^i(P - \phi_{n,j}) dP. \quad (29)$$

In turn, the outage probability and BLER for a specific layer [(30) and (32)] or service [(31) and (33)] are shown as follows:

$$\overline{\text{BLER}}_\eta^i = \frac{\sum_{n=0}^{K-1} \sum_{j=0}^{TFS_n-1} \epsilon_n \rho_{i,n} p_{n,j} \vartheta_{i,n,j}}{\sum_{s=0}^{K-1} \epsilon_s \rho_{i,s}} \quad (30)$$

$$\overline{\text{BLER}}_{\eta,n} = \frac{\sum_{i=0}^{N-1} \sum_{j=0}^{TFS_n-1} \alpha_i \rho_{i,n} p_{n,j} \vartheta_{i,n,j}}{\sum_{i=0}^{N-1} \alpha_i \rho_{i,n}} \quad (31)$$

$$\mu_\eta^i = \frac{\sum_{n=0}^{K-1} \sum_{j=0}^{TFS_n-1} \epsilon_n \rho_{i,n} p_{n,j} \lambda_{i,n,j}}{\sum_{s=0}^{K-1} \epsilon_s \rho_{i,s}} \quad (32)$$

$$\mu_{\eta,n} = \frac{\sum_{i=0}^{N-1} \sum_{j=0}^{TFS_n-1} \alpha_i \rho_{i,n} p_{n,j} \lambda_{i,n,j}}{\sum_{i=0}^{N-1} \alpha_i \rho_{i,n}}. \quad (33)$$

Notice that (30)–(33) relate the performance in terms of outage probability and BLER with a specific value of the load factor. Consequently, they allow obtaining the maximum load factor that ensures the desired performance constraints (i.e., the maximum BLER or the maximum outage probability). It is worth mentioning that this maximum load factor establishes the maximum total interference level allowed in the cell, regardless of how it is distributed between intercell or intracell interference. Therefore, although the model considers a single cell (i.e., the cell under study), the maximum obtained load factor would also be applicable in a multicell scenario (i.e., in such a case, the total maximum interference allowed in the cell would be the same, but depending on the scenario, the specific value of inter- or intracell interference could be different).

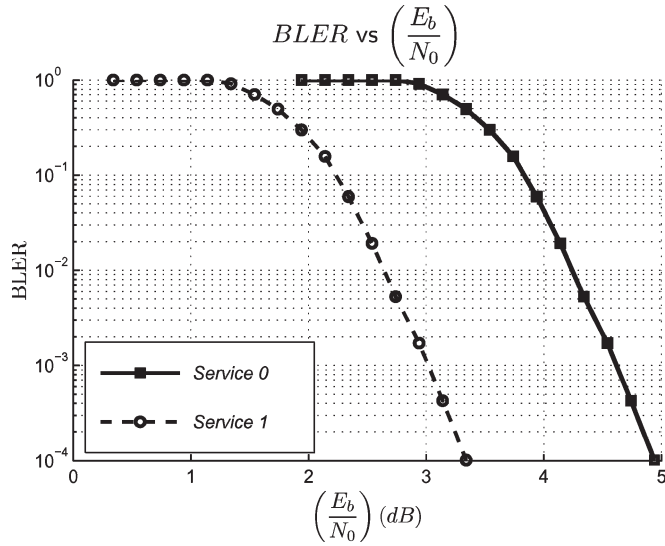
It is worth noting that all the expressions developed so far have been obtained by considering circular layers. However, the methodology would also be applicable to any other shape if propagation loss density functions are known. It is also remarkable that, for instance, it is possible to follow the same process if a measurement campaign is carried out and there is a previous knowledge of the area that is to be analyzed in terms of total propagation loss distribution. Similarly, if soft handover was to be included, this would basically modify the path loss of the users of the cell under study in (10) because it would be possible that some users located in other cells were connected to the cell under study, while the rest of the procedure would basically be the same.

#### IV. RESULTS

This section presents three relevant examples in order to provide insight on the proposed model and to illustrate its capabilities under different traffic and service distributions. In addition, the validation of the model through simulations is given. All results have been obtained considering the propagation model exposed in (6) plus a lognormal shadowing. In particular, the parameters have been set to  $Y_0 = 128.1$ ,  $\zeta = 37.6$ ,  $\sigma = 3$  dB,  $P_N = -103$  dBm,  $P_{T\max} = 21$  dBm, and  $P_{T\min} = -44$  dBm. Two services (service 0 and service 1) are considered. Service parameters are detailed in Table I. Service 0 intends to represent a conversational case-based reasoning service that generates traffic at 64 kb/s. On the other hand, service 1 intends to represent an interactive service with four possible transmission bit rates. This makes the selection of a transmission bit rate a key issue in terms of the generated interference and QoS parameters such as outage probability and BLER. The physical layer has been characterized by the function  $b_{n,j}(\cdot)$ ,

TABLE I  
 SERVICE PARAMETERS

Parameter	Service 0	Service 1
$\epsilon_n$	0.2	0.08
$TF'S_n$	1	4
$Rb_{n,0}$	64 kbps	16 kbps
$\left(\frac{E_b}{N_0}\right)_{n,0}$	4.0 dB	2.5 dB
$Rb_{n,1}$	N/A	32 kbps
$\left(\frac{E_b}{N_0}\right)_{n,1}$	N/A	2.5 dB
$Rb_{n,2}$	N/A	64 kbps
$\left(\frac{E_b}{N_0}\right)_{n,2}$	N/A	2.5 dB
$Rb_{n,3}$	N/A	128 kbps
$\left(\frac{E_b}{N_0}\right)_{n,3}$	N/A	2.5 dB


 Fig. 2. Service BLER versus  $E_b/N_0$ .

as shown in Fig. 2. Notice that  $\gamma_{n,j}(\cdot)$  is obtained by solving  $\gamma_{n,j}(P) = b_{n,j}(g(P))$  [see (26)]. Curves shown in Fig. 2 have been obtained through a link level simulator that includes the 1500-Hz closed-loop power control, 1/3 turbo coding effect, and channel impulse response estimation.

#### A. Case Study 1

In the first analysis approach, a nonuniform traffic spatial distribution is considered in a single service scenario including only service 0. The scenario is composed of two traffic layers ( $L_0$  and  $L_1$ ), as shown in Fig. 3. Case study parameters are listed in Table II. It represents a hotspot region where there is a higher concentration of traffic.

Fig. 4(a) and (b) plots the pdf of the required transmitted power for different  $D_1$  values when  $\alpha_1$  is varied. In both figures, a load factor  $\eta = 0.8$  is considered.  $\alpha_1$  is varied from 0 to 1. When  $\alpha_1 = 0$ , there are no users within layer  $L_1$ , and thus, users are uniformly distributed all around a cell of radius  $r_0$ . On the other hand, when  $\alpha_1 = 1$  all users are spread over layer  $L_1$ , whose center is placed 150 m [Fig. 4(a)] or 950 m [Fig. 4(b)] from the cell site. The distribution of the total required transmitted power is the weighted sum of the distributions of layers  $L_0$  and  $L_1$  components. Thus, for small

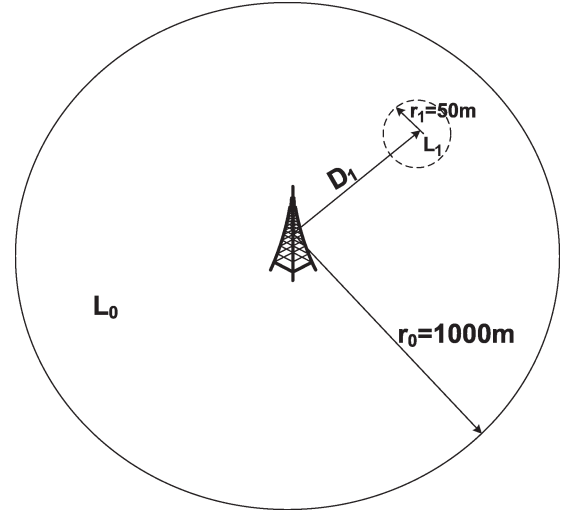


Fig. 3. Case study 1 layout.

 TABLE II  
 CASE STUDY 1 PARAMETERS

Parameter	Value
$r_0$	1000 m
$D_0$	0 m
$\alpha_0$	$1-\alpha_1$
$r_1$	50 m
$D_1$	[50 m, 950 m]
$\alpha_1$	[0.0, 1.0]

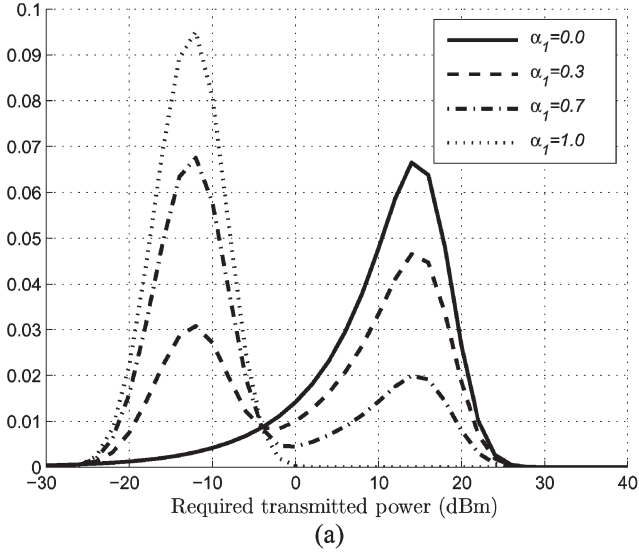
$D_1$  values two peaks arise in the total pdf (one for each layer). As distance  $D_1$  is increased, the peak associated with layer  $L_1$  is moved to higher values. For high  $D_1$  values, contributions of the two layers result in a single peak [Fig. 4(b)]. For instance, when  $D_1 = 150$  m, peaks are centered at about  $-12$  and  $15$  dBm. If  $D_1 = 950$  m, a single peak is centered at about  $15$  dBm. It is also remarkable that for high  $D_1$ , the single pdf peak narrows as  $\alpha_1$  increases because users are more concentrated spatially.

The validation of the model has been carried out through simulations. Fig. 4(c) plots a comparison between required transmitted power pdfs obtained theoretically and with simulations. It is shown that differences are minor, and the model fits the results obtained by simulation.

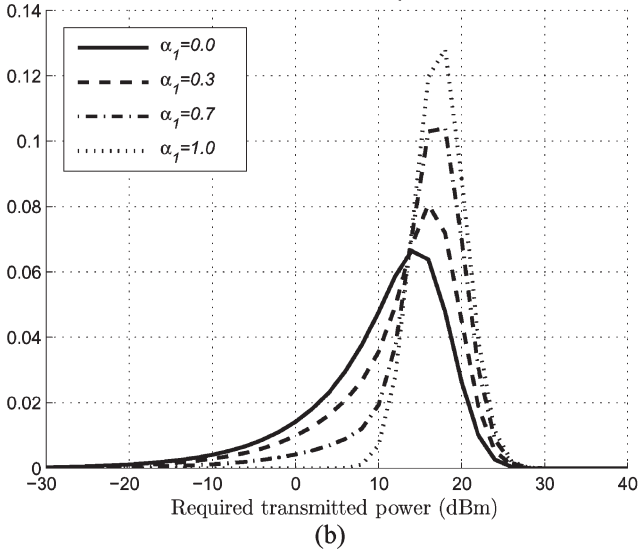
As transmitted power is upper bounded by  $P_{Tmax}$ , the performance of the system is tightly coupled with values over this threshold (in this particular case, 21 dBm). Fig. 5(a) and (b) shows the outage associated to case study 1. In these figures, two different situations are examined, namely 1)  $L_1$  is close to the base station [Fig. 5(a)], and 2)  $L_1$  is far from the base station [Fig. 5(b)]. Taking into account that outage will be lower if users are concentrated in the vicinity of the base station, in Fig. 5(a), high  $\alpha_1$  implies low outage probability, whereas in Fig. 5(b), high  $\alpha_1$  implies high outage.

The same interpretation stated for outage probability is valid for BLER. Fig. 5(c) and (d) shows the average BLER in the same conditions in which outage was obtained. The target BLER is set to 5%, which is associated with an  $E_b/N_0 = 4$  dB. All results have been obtained analytically.

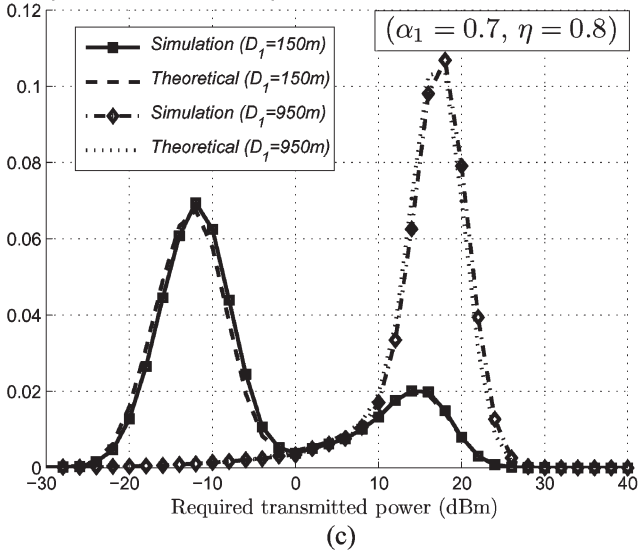
Required transmitted power pdf ( $D_1 = 150m, \eta = 0.8$ )



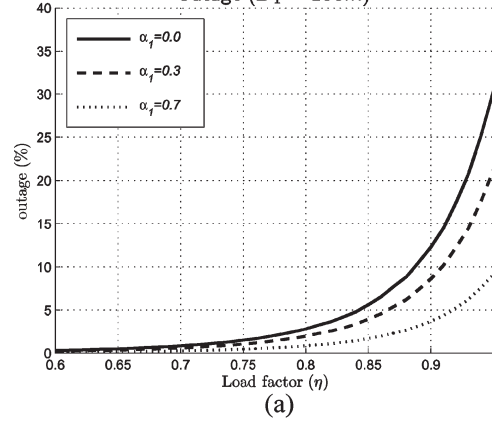
Required transmitted power pdf ( $D_1 = 950m, \eta = 0.8$ )



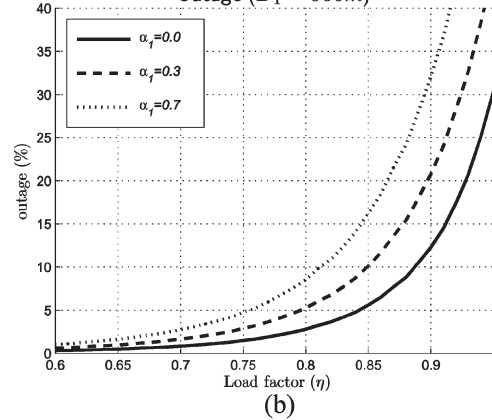
Required transmitted power: theoretical vs. simulation



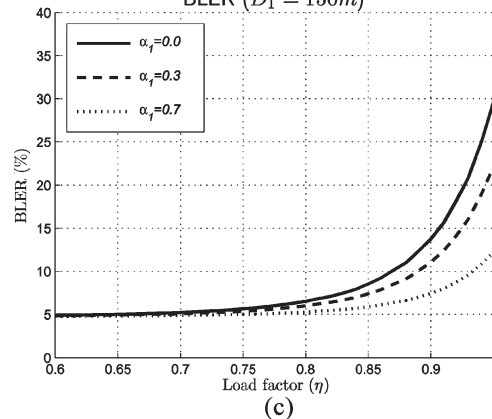
Outage ( $D_1 = 150m$ )



Outage ( $D_1 = 950m$ )



BLER ( $D_1 = 150m$ )



BLER ( $D_1 = 950m$ )

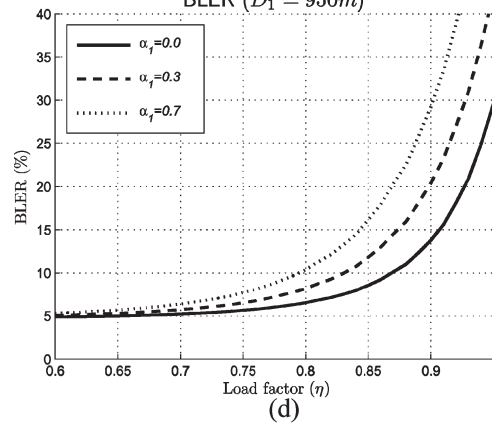


Fig. 4. Required transmitted power pdf (a) with  $D_1 = 150m$ , (b) with  $D_1 = 950m$ , and (c) validation of the model.

Fig. 5. Outage [with (a)  $D_1 = 150m$  and (b)  $D_1 = 950m$ ] and BLER [with (c)  $D_1 = 150m$  and (d)  $D_1 = 950m$ ].

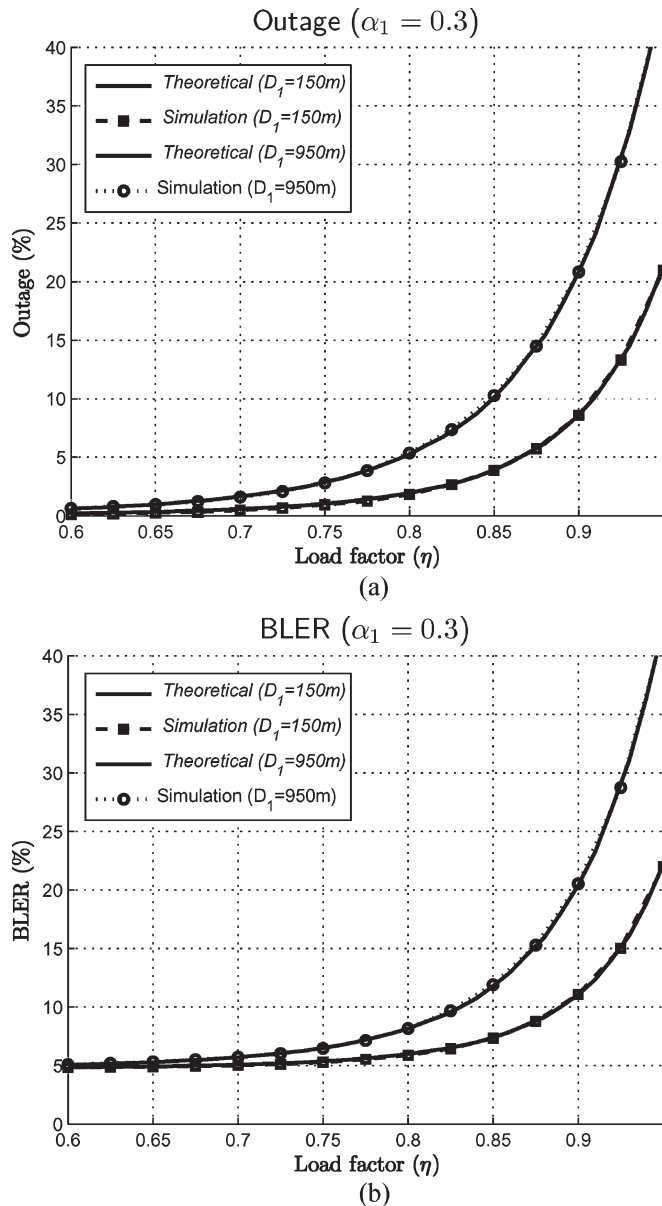


Fig. 6. Validation of (a) outage and (b) BLER with  $\alpha_1 = 0.3$ .

In order to validate such results, Fig. 6(a) and (b) compares the results obtained theoretically and through simulations. In both figures,  $\alpha_1$  has been set to 0.3 and  $D_1$  is set to 150 or 950 m. Differences observed between simulation and theory are negligible.

If a maximum outage probability threshold  $\mu_{\text{Threshold}}$  is set, it is possible to determine the maximum cell capacity, which is given by the maximum load factor level  $\eta_{\text{max}}$  that corresponds to this outage threshold. This maximum load factor level could be the input to the different uplink RRM algorithms, e.g., admission or load control. Fig. 7(a) and (b) shows the maximum allowable load factor level for an outage probability threshold  $\mu_{\text{Threshold}} = 0.5\%$ . The former shows the evolution of the maximum allowable load when  $D_1$  is varied. The latter depicts the same evolution as a function of  $\alpha_1$ . It is observed in Fig. 7(a) that high  $\alpha_1$  values lead to higher capacity (high load) when  $D_1$  is small because most users are concentrated

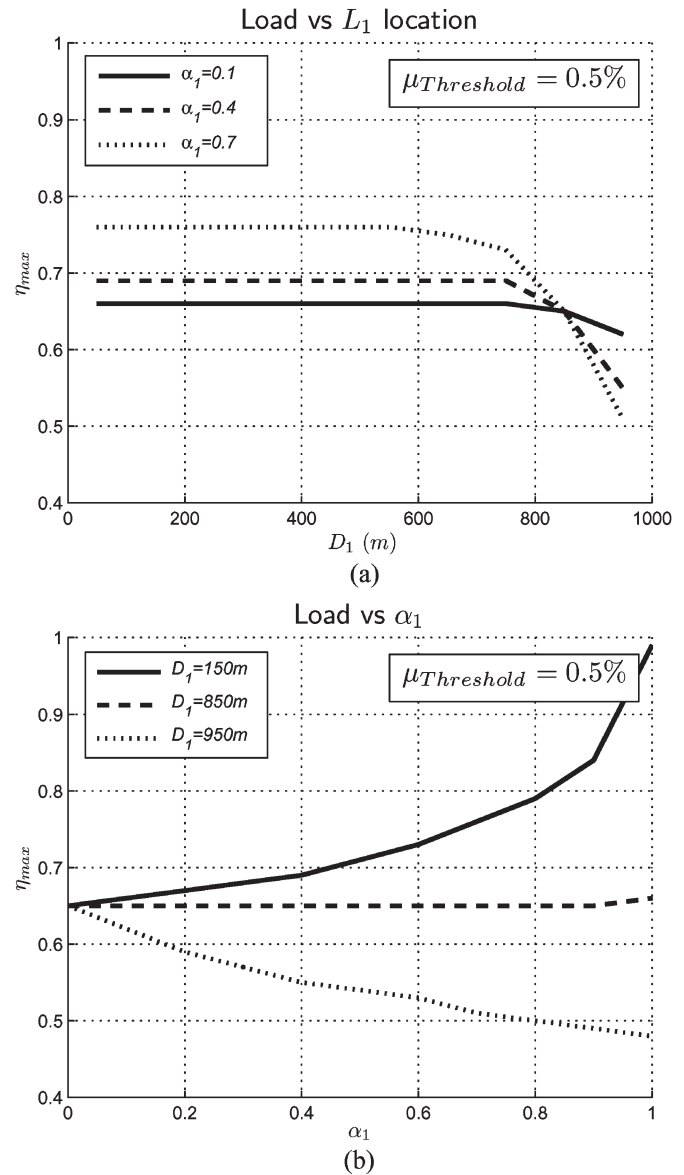


Fig. 7. Load factor corresponding to  $\mu_{\text{Threshold}} = 0.5\%$  as a function of (a)  $D_1$  and (b)  $\alpha_1$ .

close to the base station. On the other hand, it implies a decrease of capacity as  $D_1$  increases. Fig. 7(b) plots the same effect by varying  $\alpha_1$ .

Note that in Fig. 7(a) outage level matches up for all cases ( $\alpha_1 = 0.1$ ,  $\alpha_1 = 0.4$ , and  $\alpha_1 = 0.7$ ) when  $D_1 = 850$  m, and in fact, it could be extended to any  $\alpha_1$  value. This effect is due to the total loss pdfs of each layer ( $L_0$  and  $L_1$ ). When  $r_0 = 1000$  m and  $D_1 = 850$  m, both pdfs are centered at the same range of values. Then, the final total loss pdf does not vary significantly for different  $\alpha_0$  and  $\alpha_1$  values, assuming that only two layers are considered. As the outage probability is computed from the final pdf, the outage is similar for all  $\alpha_1$  values.

Three different trends may be observed in Fig. 7(b). When  $D_1 = 150$  m, the maximum load factor  $\eta_{\text{max}}$  increases with  $\alpha_1$ ; when  $D_1 = 850$  m, the maximum load factor remains constant; finally, when  $D_1 = 950$  m, the maximum load decreases as



TABLE III  
CASE STUDY 2 PARAMETERS

Parameter	Value
$r_0$	1000 m
$D_0$	0 m
$\alpha_0$	0.5
$r_1$	50 m
$D_1$	150 m
$\alpha_1$	0.5

TABLE IV  
SERVICE DISTRIBUTION

$\rho_{i,n}$	$S_0$	$S_1$	$S_2$
$\rho_{0,0}$	1.0	0.5	0.0
$\rho_{0,1}$	0.0	0.5	1.0
$\rho_{1,0}$	0.0	0.5	1.0
$\rho_{1,1}$	1.0	0.5	0.0

TABLE V  
SERVICE 1 TF SELECTION PROBABILITIES

$p_{n,j}$	$MAC_0$	$MAC_1$
$p_{1,0}$	0.0	0.5
$p_{1,1}$	0.0	0.5
$p_{1,2}$	0.0	0.0
$p_{1,3}$	1.0	0.0

$\alpha_1$  increases. In general, high users' concentrations in  $L_1$  are beneficial for low  $D_1$  scenarios, whereas low concentrations in  $L_1$  are beneficial for high  $D_1$ .

B. Case Study 2

The scenario is composed of two traffic layers ( $L_0$  and  $L_1$  as in case study 1) and two services (service 0 and service 1). Scenario parameters are detailed in Tables III and IV. Table IV presents the distribution of users among layers, with three different situations, which are denoted as  $S_0$ ,  $S_1$ , and  $S_2$ . In the first case ( $S_0$ ), all users with service 0 are scattered within layer  $L_0$ , whereas all users with service 1 remain within layer  $L_1$ . In case  $S_1$ , there is the same number of users with service 0 and service 1 in both layers (i.e., there is no service nonuniformity but traffic spatial nonuniformity). Finally, in case  $S_2$  all users with service 0 are scattered in  $L_1$  and users with service 1 in  $L_0$ .

The considered scenario models an urban environment where specific areas present traffic spatial and/or service nonuniformities (e.g., bus stations). Service 1 has several possible bit rates, and thus, there is a probability associated to each of the possible TFs. These probabilities depend on the user equipment MAC algorithms for TF selection [4]. As these MAC algorithms are not the purpose of this paper, two different sets of TF selection probabilities are considered and assumed to be the result of applying two algorithms that are labeled as  $MAC_0$  and  $MAC_1$  whose probabilities are detailed in Table V. Fig. 3 has been obtained for a definite coding scheme used in both  $MAC_0$  and  $MAC_1$  algorithms. Therefore, the curve plotted in Fig. 3 is used for both MAC algorithms. Whereas  $MAC_1$  algorithm tends to select low bit transmission rates,  $MAC_0$  selects the highest bit transmission rate, thus representing two different QoS profiles.

Let us consider the  $MAC_0$  algorithm. Fig. 8(a) shows the required transmitted power pdf with  $\eta = 0.8$ . In all cases

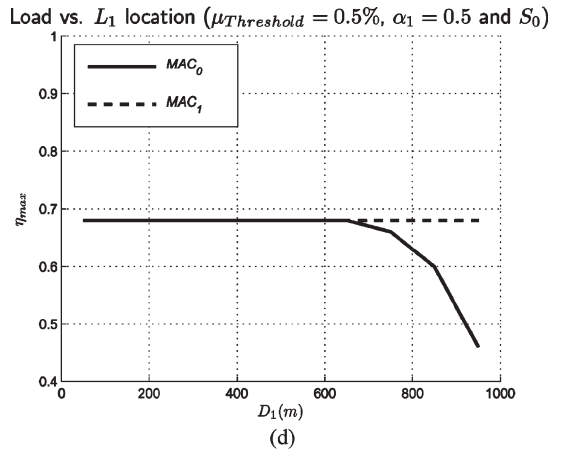
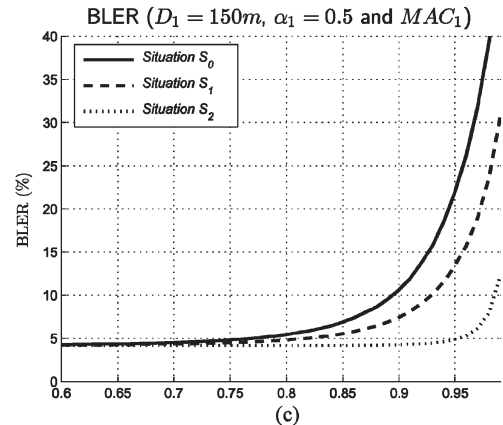
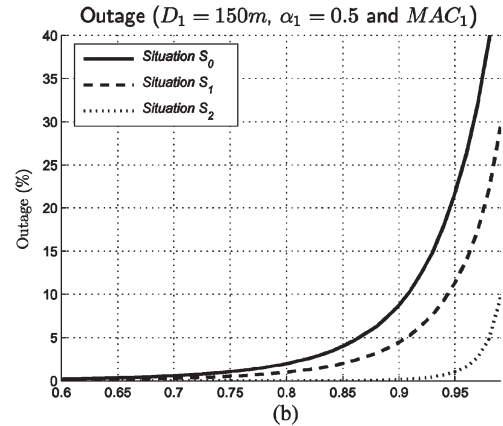
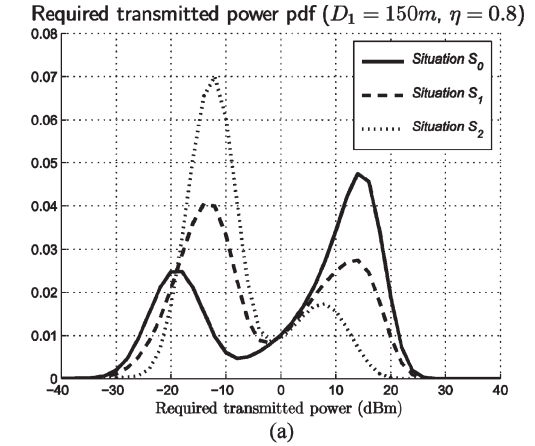


Fig. 8. (a) Required transmitted power pdf. (b) Outage. (c) BLER. (d) Load factor corresponding to  $\mu_{Threshold} = 0.5\%$  for  $MAC_0$  and  $MAC_1$ .

$(S_0, S_1, S_2)$ , two peaks can be observed. Each peak is associated to one of the layers (i.e., the one appearing at low power values is associated with  $L_1$  and the one appearing at high power values with  $L_0$ ). When different services with different activity factors ( $\epsilon_n$ ) are mixed, the ratio between activity factors determines the proportion of users of each service. Therefore, the probability of transmitting high power levels (associated to users in layer  $L_0$ ) is higher in  $S_0$  (all users belonging to service 0 are within layer  $L_0$ ) than in the others. On the other hand, when all users with service 0 are located in layer  $L_1$  (case  $S_2$ ), low transmitted power levels are more likely to occur. Finally, case  $S_1$  presents a more balanced situation between users in each layer because, although users with service 0 have a bigger impact on transmitted power, the proportion of users with each service is maintained equal in both layers.

The previous explanation can be generalized to any load factor level. However, the higher the load factor is, the higher the required power level will be, thus increasing the outage. This effect is shown in Fig. 8(b), which plots the performance of the system in terms of outage as a function of load factor. BLER is shown in Fig. 8(c). The target BLER is set to 5% for service 0 and to 2.5% for service 1.

The impact of the two MAC algorithms and their associated TF probabilities in terms of system capacity is analyzed in Fig. 8(d). The outage threshold is set to 0.5% in  $S_0$ . Notice that differences arise for high  $D_1$  values. In fact, results obtained with different MAC strategies are more significant as differences in used  $(E_b/N_0) \cdot R_b$  increase. Fig. 8(d) plots a comparison between the two algorithms detailed in Table V as a function of  $D_1$ . In this case, the capacity is reduced when high transmission bit rates are more likely to be selected ( $MAC_0$ ). The capacity in terms of load factor is maintained for the whole range of  $D_1$  with  $MAC_1$  because selected transmission bit rates are low and do not imply an increase of the outage even for high  $D_1$ . When  $MAC_0$  is used, the selection of high transmission bit rate causes an increase on outage, and thus, the capacity is reduced.

### C. Case Study 3

In this section, an example of a more complex scenario is analyzed (Fig. 9). The scenario is now composed of four traffic layers and the same two services used in case study 2 (Table I). Parameters are shown in Tables VI and VII. In the first layer ( $L_0$ ), there are only users with service 0. Three nonuniformities are placed within the cell (layers  $L_1$ ,  $L_2$ , and  $L_3$ ), in which only users with service 1 are considered. Service 1 uses the  $MAC_1$  algorithm detailed in Table V.

In previous case studies, the interpretation of the concrete results obtained with the derived expressions was not difficult to be carried out due to the reduced number of services and layers under study. However, as the number of layers and/or services increases, the interpretation becomes less obvious. Therefore, in such case studies, the proposed analytical model becomes a suitable tool to evaluate the cell capacity. Case study 3 intends to show an example of these complex scenarios. Fig. 10(a) plots the required transmitted power pdf where all effects of traffic spatial and service distributions are mixed up.

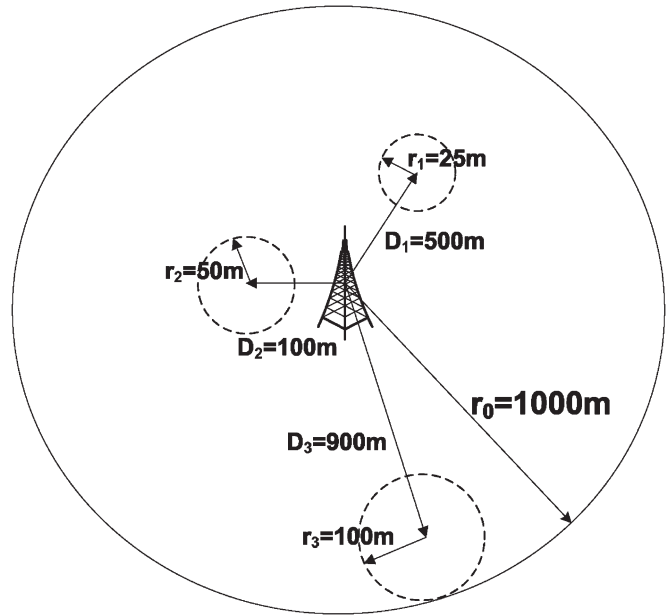


Fig. 9. Case study 3 layout.

TABLE VI  
CASE STUDY 3 PARAMETERS

Parameter	Value
$r_0$	1000 m
$D_0$	0 m
$\alpha_0$	0.4
$r_1$	25 m
$D_1$	500 m
$\alpha_1$	0.2
$r_2$	50 m
$D_2$	100 m
$\alpha_2$	0.2
$r_3$	100 m
$D_3$	900 m
$\alpha_3$	0.2

TABLE VII  
SERVICE DISTRIBUTION

Parameter	Value
$\rho_{0,0}$	1.0
$\rho_{0,1}$	0.0
$\rho_{1,0}$	0.0
$\rho_{1,1}$	1.0
$\rho_{2,0}$	0.0
$\rho_{2,1}$	1.0
$\rho_{3,0}$	0.0
$\rho_{3,1}$	1.0

Fig. 10(b) shows the total, service 0, and service 1 outage probabilities. Users with service 1 present the best performance in terms of outage probability due to the distance at which  $L_1$  and  $L_2$  layers are located. Sixty percent of the users use service 1. Two out of three users with service 1 belong to layer  $L_1$  or  $L_2$ . As  $L_1$  and  $L_2$  have their center at 500 and 100 m, respectively, and their radii are 25 and 50 m, there are no users of layers  $L_1$  and  $L_2$  at a distance higher than 525 m. Forty percent of the users use service 0 and are spread in layer  $L_0$ , whose radius is 1000 m. Therefore, outage of users with service 1 is lower than that of users with service 0. Finally, the total outage is the sum of the contribution of both services.

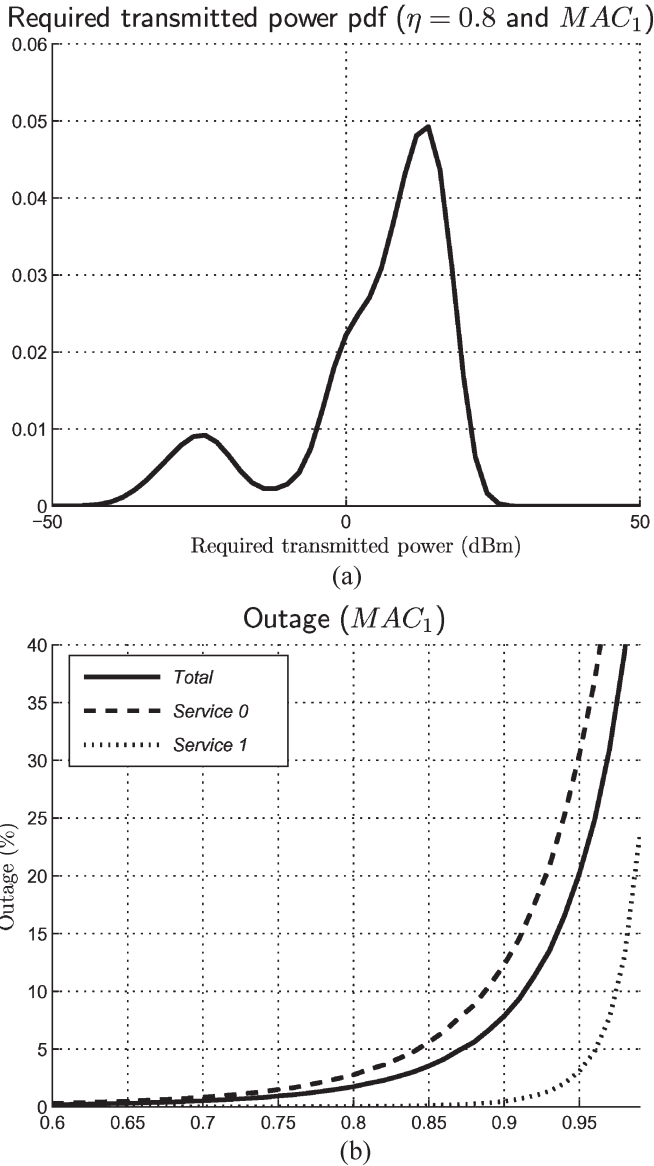


Fig. 10. (a) Required transmitted power. (b) Total, service 0, and service 1 outage probabilities for case study 3.

Despite having 60% of the users using service 1 and only 40% using service 0, the total outage curve is close to the curve of service 0 because of the higher activity factor of service 0. In particular, if the threshold outage  $\eta_{\text{Threshold}}$  is set to 0.5%, the maximum load factor for case study 3 is 0.69.

### V. CONCLUSION

An analytical model has been proposed for the reverse link of a WCDMA system with traffic nonuniform distribution and supporting multiple services. The main contribution of the article lies on the generalization of hotspots in terms of spatial traffic nonuniformities as well as in terms of service distribution nonuniformities. QoS parameters, such as BLER and outage probability, have been characterized for a circular cell with circular nonuniformities or hotspots, though the description of the process may also be applied to any known scenario by properly replacing the path loss pdfs. Therefore, the proposed

model allows a detailed study of the impact caused by the mix of different services in a nonuniformly distributed scenario and the computation of the maximum capacity (i.e., the maximum uplink load factor that can be supported for specific QoS constraints). A proper maximum load selection becomes a key point to maximize the capacity, whereas average outage or BLER probabilities are maintained below a design value. An excessively high load threshold could degrade the performance of the system. On the other hand, too low loads would not be efficient in terms of available resources usage. Some case studies have been analyzed, and the maximum load factor has been obtained.

Different conclusions have been highlighted in the analyzed case studies. In particular, it has been observed that the maximum admissible load factor  $\eta_{\text{max}}$  is dependent on the traffic distribution, which is characterized in this paper through the number of traffic layers and the location and the proportion of traffic in each layer. The variability of  $\eta_{\text{max}}$  is shown to be higher when hotspots are located either in the vicinity of the base station or close to the cell border, whereas it presents a more stable behavior for intermediate locations. Likewise, MAC algorithms have an important impact on the capacity of the system, and those MAC algorithms that use high transmission bit rates lead to lower maximum load factors.

Finally, the complexity to determine the maximum load factor increases as the number of layers and/or services is increased. The potentials of the model are stressed in these scenarios, where the prediction of the load values is to be used as RRM algorithm input. In particular, one of the most important RRM strategies is the CAC. This procedure is usually carried out by setting a maximum allowable load factor  $\eta_{\text{max}}$ , which is used to decide whether a new connection should be admitted or rejected. Usually, this maximum load is determined regardless of the traffic distribution. It has been demonstrated throughout this paper that in scenarios with traffic nonhomogeneities this value varies. The maximum allowable load factor can be easily obtained in simple scenarios. However, as the complexity of the scenario increases, it becomes more and more difficult to predict  $\eta_{\text{max}}$ ; thus, the proposed methodology becomes a useful tool.

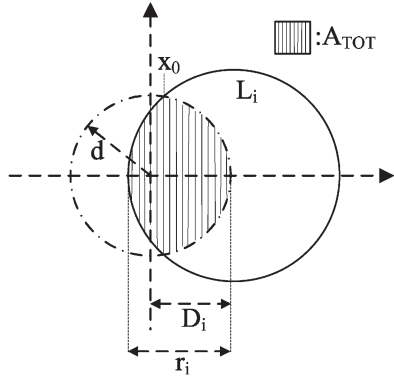
### APPENDIX

In the following, the pdf of the path loss of a traffic layer, as defined in Section II, is computed. As shown in (6), the path loss is dependent on the distance to the base station ( $d$ ). Then, the first step to obtain the pdf of the path loss ( $f_Y(y)$ ) is the calculation of the distance pdf ( $f_D(d)$ ) of a generic round-shaped layer  $i$  with radius  $r_i$  (in meters) and whose center is placed  $D_i$  meters from the coordinate origin. The distance probability distribution function presents two different cases.

*Case 1— $D_i < r_i$ :* If  $D_i < r_i$ , the distance probability distribution function of layer  $L_i$  is defined as

$$F_D^i(d) = \text{prob}(d \leq D) = \frac{A_{\text{TOT}}}{\pi r_i^2} \tag{34}$$

where  $A_{\text{TOT}}$  is the layer  $L_i$  area intersected with a circle of radius  $d$ , as shown in Fig. 11. Notice that  $A_{\text{TOT}} = \pi d^2$  when


 Fig. 11. Scheme of the situation when  $D_i < r_i$  and  $d > r_i - D_i$ .

$d \leq r_i - D_i$  (i.e., the circle of radius  $d$  is inside the circle of radius  $r_i$ ). However, when  $d > r_i - D_i$ ,  $A_{TOT}$  is given as follows (see Fig. 11):

$$A_{TOT} = 2 \int_{D_i - r_i}^{x_0} \sqrt{(r_i^2 - (x - D_i)^2)} dx + 2 \int_{x_0}^d \sqrt{(d^2 - x^2)} dx \quad (35)$$

where

$$x_0 = \frac{D_i^2 + d^2 - r_i^2}{2D_i} \quad (36)$$

as shown in Fig. 11. By computation, we have

$$\begin{aligned} A_{TOT} = & \frac{\pi}{2} (d^2 + r_i^2) - d \arcsin \left( \frac{D_i + d^2 - r_i^2}{2dD_i} \right) \\ & + r_i^2 \arcsin \left( \frac{d^2 - D_i - r_i^2}{2r_iD_i} \right) \\ & - \frac{1}{2} \sqrt{(d^2 - (D_i - r_i)^2) ((D_i - r_i)^2 - d^2)}. \end{aligned} \quad (37)$$

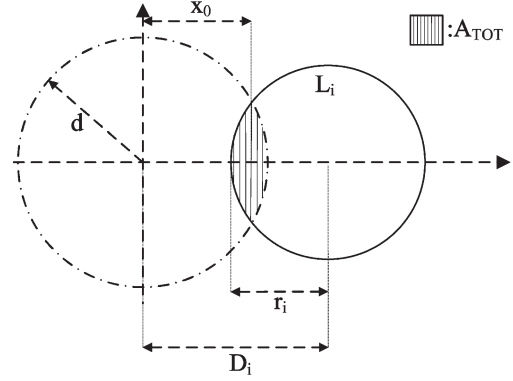
Therefore, substituting (37) in (34) yields

$$f_D^i(d) = \frac{d^2}{r_i^2} \quad (38)$$

if  $0 \leq d \leq (r_i - D_i)$ , or

$$\begin{aligned} f_D^i(d) = & \frac{1}{\pi r_i^2} \left[ \frac{\pi}{2} (d^2 + r_i^2) - d^2 \arcsin \left( \frac{D_i^2 + d^2 - r_i^2}{2D_i d} \right) \right. \\ & + r_i^2 \arcsin \left( \frac{d^2 - D_i^2 - r_i^2}{2D_i r_i} \right) \\ & \left. - \frac{1}{2} \sqrt{(d^2 - (D_i - r_i)^2) ((D_i + r_i)^2 - d^2)} \right] \end{aligned} \quad (39)$$

if  $(r_i - D_i) \leq d \leq (D_i + r_i)$ .


 Fig. 12. Location of the generic layer when  $D_i \geq r_i$ .

The pdf is obtained by deriving (38) and (39) as a function of distance  $d$ , i.e.,

$$f_D^i(d) = \frac{2d}{r_i^2} \quad (40)$$

in the range  $0 \leq d \leq (r_i - D_i)$ , and

$$f_D^i(d) = \frac{1}{\pi r_i^2} \left[ \pi d - 2d \arcsin \left( \frac{D_i^2 + d^2 - r_i^2}{2D_i d} \right) \right] \quad (41)$$

in the range  $(r_i - D_i) \leq d \leq (D_i + r_i)$ .

*Case 2— $D_i \geq r_i$ :* If  $D_i \geq r_i$  (Fig. 12),  $A_{TOT}$  is analogous to the case obtained in (37), and thus, the probability distribution function is expressed as

$$\begin{aligned} F_D^i(d) = & \frac{1}{\pi r_i^2} \left[ \frac{\pi}{2} (d^2 + r_i^2) - d^2 \arcsin \left( \frac{D_i^2 + d^2 - r_i^2}{2D_i d} \right) \right. \\ & + r_i^2 \arcsin \left( \frac{d^2 - D_i^2 - r_i^2}{2D_i r_i} \right) \\ & \left. - \frac{1}{2} \sqrt{(d^2 - (D_i - r_i)^2) ((D_i + r_i)^2 - d^2)} \right] \end{aligned} \quad (42)$$

in the range  $(D_i - r_i) \leq d \leq (D_i + r_i)$ .

Therefore, the pdf is given as follows:

$$f_D^i(d) = \frac{1}{\pi r_i^2} \left[ \pi d - 2d \arcsin \left( \frac{D_i^2 + d^2 - r_i^2}{2D_i d} \right) \right] \quad (43)$$

in the range  $(D_i - r_i) \leq d \leq (D_i + r_i)$ .

Once distance distribution is obtained, path loss distribution may be calculated through the relationship existing between distance  $d$  and path loss  $Y$  extracted from (6). In linear units, distance is expressed as follows:

$$d = 10^{\frac{Y - Y_0}{\zeta}} = g^{-1}(Y). \quad (44)$$

In this context,  $f_Y(y)$  can be obtained from  $f_D(d)$ . Since  $d$  is a continuous random variable and  $g^{-1}(Y)$  is a differentiable

function,  $f_Y(y)$  is given by

$$f_Y^i(y) = f_D^i(g^{-1}(y)) \left| \frac{dg^{-1}(y)}{dy} \right|. \quad (45)$$

Then, introducing (40), (41), (43), and (44) in (45), it finally yields the pdf of the propagation loss given in (7)–(9).

#### ACKNOWLEDGMENT

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